

ORTHOGONAL SYSTEMS OF FUNCTIONS (chapter 2)

We start from definition of scalar product  $(f, g)$

$$(f, g) \stackrel{(*)}{=} \int_a^b f(x) \overline{g(x)} w(x) dx$$

There are several possible scalar products, depending on the choice of the interval  $[a; b]$ , and the weight function  $w(x) \geq 0$  (a very simple choice of weight is  $w(x) = 1$  or  $w(x) = c$  a constant)

In classical Fourier Series  $[a; b] = [-\pi; \pi]$ , but it could be any interval of length  $T > 0$  (period of  $f$ ). The interval could also be  $[0; +\infty)$  a half-line or  $(-\infty; +\infty)$  a whole line.

Chosen a specific scalar product

Def. we say that  $f \perp g$  with respect to the scalar product  $(*)$  if and only if  $\int_a^b f(x) \overline{g(x)} w(x) dx = 0$

(the integral = 0 because of "cancellation effect" of positive and negative areas)

Remark  $(f, f) = \int_a^b |f(x)|^2 w(x) dx \geq 0$

$$\|f\|^2 \text{ on } L^2([a, b], w)$$

Consider a set of ( $\infty$  many) functions  $\phi_n(x)$  with  $x \in (a, b)$ , and  $n \in I$  (set of indices, usually  $\mathbb{N}^+, \mathbb{N}^0, \mathbb{Z}$ ).

This is called orthogonal system if:

$$(1) \phi_n \in L^2([a, b], w) \quad \forall n \in I$$

$$(2) \|\phi_n\|_2 > 0 \quad \forall n \in I \quad \text{where}$$

(the norm must be positive)

$$\|\phi_n\|_2 = \|\phi_n\| = \left( \int_a^b |\phi_n(x)|^2 w(x) dx \right)^{1/2}$$

$$(3) (\phi_n, \phi_m) = 0 \quad \text{if } n \neq m$$

(these functions are  $\perp$  to each other)

### Example (important)

$$\phi_m(x) = e^{imx} = \cos mx + i \sin mx$$

$$I = \mathbb{Z} = \{ \dots, -3, -2, -1, 0, 1, 2, 3, \dots \}$$

$[a; b] = [-\pi; \pi]$  (actually any interval of length  $2\pi$  is OK, e.g.  $[0; 2\pi]$  or  $[-\frac{\pi}{4}; \frac{7}{4}\pi]$ )

$$w(x) = 1 \quad L^2([-\pi; \pi], dx)$$

$$(1) \int_{-\pi}^{\pi} |e^{imx}|^2 dx = \int_{-\pi}^{\pi} 1 dx = 2\pi \Rightarrow \|\phi_m\|_2 = \sqrt{2\pi}$$

$$|e^{imx}|^2 = |\cos mx + i \sin mx|^2 = \cos^2 mx + \sin^2 mx = 1$$

$$(2) \|\phi_m\|_2 = \sqrt{2\pi}$$

N.B.: (1) is almost trivial in this case, because  $\phi_m$  itself is bounded and  $\int_a^b |\phi_m(x)|^2 dx = \infty$  impossible, BUT for general orthogonal systems the  $\phi_m(x)$  could be unbounded; as long as  $\|\phi_m\| < \infty$  FINITE.

$$(3) (\phi_n, \phi_m) = \int_{-\pi}^{\pi} \underbrace{e^{inx}}_{\phi_n} \cdot \underbrace{e^{-imx}}_{\overline{\phi_m}} dx = \int_{-\pi}^{\pi} e^{i(n-m)x} dx = \begin{matrix} \text{with} \\ n-m \neq \\ \text{INTEGER} \end{matrix} \\ = \left[ \frac{e^{i(n-m)x}}{i(n-m)} \right]_{-\pi}^{\pi} = 0 \quad \begin{matrix} \text{REMEMBER} \\ e^{i\pi} = e^{-i\pi} \end{matrix} \quad \begin{matrix} \text{|||} \\ 0 \text{ because of} \\ \text{cancellation} \end{matrix}$$

Def.: suppose  $f(x) \in L^1([a, b], w)$  and also  $f(x) \overline{\phi_m(x)} \in L^1([a, b], w)$  for  $m \in I$ , then

$$c_m = \frac{1}{\|\phi_m\|_2^2} \cdot (f, \phi_m)$$

are the generalized Fourier coefficients of  $f(x)$  w.r.t. the orthogonal system  $\{\phi_m\}$ .

N.B. There are many orthogonal systems besides the classical  $e^{inx}$ ...

N.B. When  $\phi_m(x) = e^{inx}$ ,  $I = \mathbb{Z}$ , then  $c_n$  are the classical Fourier coeff's of the  $2\pi$ -periodic function  $f$ :

$$c_m = \frac{1}{2\pi} \int_{-\pi}^{\pi} f(x) e^{-imx} dx$$

$\| \phi_m \|_2^2 \nearrow \underbrace{\hspace{10em}}_{(f, e^{imx})}$

N.B. if instead of  $2\pi$ -periodic functions we want to consider  $T$ -periodic functions, then  $\phi_m(x) = e^{i \frac{2\pi}{T} mx}$ ,  $[a, b] = [-\frac{T}{2}, \frac{T}{2}]$ , or any interval of length  $T$ , like e.g.  $[0, T]$ .

In the  $T$ -periodic case, our scalar product will be: ( $w=1$ )

$$(f, g) = \int_{-T/2}^{T/2} f(x) \overline{g(x)} dx$$

$$c_m = \frac{1}{T} \int_{-T/2}^{T/2} f(x) e^{-i \frac{2\pi}{T} mx} dx$$

Remark F.S. are well adapted both to the case of  $T$ -periodic functions and the case of functions that are only defined in  $[a; b]$ , with  $(b-a = T)$ , because we can always assume that a function defined only on  $[a; b]$  is periodic outside.

Guideline (Road Map) to apply Fourier Theory to PDE's:

- (1) Domain has a side which is a finite interval  $\Rightarrow$  Fourier Series (classical or generalized)
- (2) Domain has a full line in the boundary  $\Rightarrow$  Fourier Transform
- (3) Domain has a half line on the boundary  $\Rightarrow$  Laplace Transform

The expression  $\sum_{n \in I} c_n \phi_n(x)$  is called the expansion of  $f(x)$

as Fourier Series (Generalized) w.r.t. the orthogonal system  $\{ \phi_n \}$ .

This F.S. converges to  $f(x)$  at least in the  $L^2$  sense

e.g.  $\lim_{N \rightarrow \infty} \|f - \sum_{n=-N}^N c_n \phi_n(x)\|_2 = 0$

Then, if  $f$  is "good enough" (for example  $f \in C^1[a; b]$ ), then we have other kinds of convergence ( $C^1 \Rightarrow$  pointwise, uniform...)

N.B. If  $I = \mathbb{Z}$  the partial sums are from  $-N$  to  $N$  (balanced, symmetric, finite sums). If  $I = \{0, 1, 2, 3, \dots\}$ , then the partial sums are from  $0$  to  $N$ , etc.

Def. An orthogonal system is called ortho-normal if

$$\|\phi_n\|_2 = 1 \text{ for } n \in I$$

Any orthogonal system can be normalized, just multiplying each  $\phi_n$  by a suitable constant.

Theorem (minimizing property of Fourier Coefficients)

$$f(x) \sim \sum_{n \in I} c_n \phi_n(x) \quad f \in L^2([a; b]) \quad c_n = \frac{1}{\|\phi_n\|_2} \cdot (f, \phi_n)$$

$$\|f(x) - \sum_{j=-N}^N c_j \phi_j(x)\|_2 \leq \|f(x) - \sum_{j=-N}^N d_j \phi_j(x)\|_2$$

and the  $=$  holds  $\iff d_j = c_j$  (Fourier coeff's  $\forall j$ )

We could say that the Fourier partial sums ( $\sum_{-N}^N \dots$ ) give us the best  $L^2$  projection on a vector sub-space of  $L^2$  of finite dimension  $(2N+1)$  (Proof later...)

Corollary  $\|f(x) - S_N(x)\|_2^2 = \|f\|_2^2 - \sum_N |c_j|^2 \|\phi_j\|_2^2$

Corollary of the Corollary (Bessel's inequality):

$$\sum_{j \in I} |c_j|^2 \|\phi_j\|_2^2 \leq \|f\|_2^2$$

because  $\|f(x) - S_N(x)\|_2$  is positive, so  $f(x) > S_N(x)$

Def. We say that  $\{\phi_n\}$  orthogonal system has the Parseval Property if, for  $\forall f \in L^2([a; b])$ , we have the equality in the Bessel's inequality:

$$\sum_{j \in I} |c_j|^2 \|\phi_j\|_2^2 = \|f\|_2^2 \quad (\infty\text{-dimensional Pythagorean theorem})$$

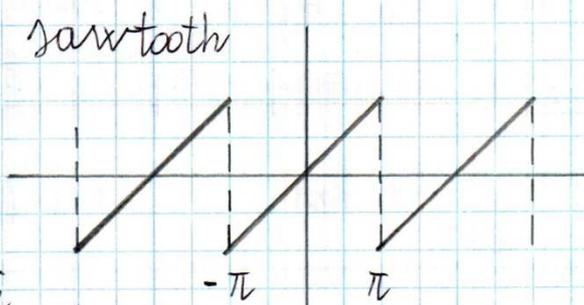
(\*)

Def. An orthogonal system  $\{\phi_n\}$  is complete if no other function can be added to it, preserving properties (1) (2) (3) (in particular prop. (3), orthogonality).

Theorem: if an orthogonal system has the Parseval Property (\*), then it is complete.

Example:

$$f(x) \begin{cases} x \text{ for } x \in (-\pi; \pi) \\ 2\pi\text{-periodic} \end{cases}$$



$$f(x) = \sum_{k=-\infty}^{+\infty} c_k e^{ikx} \quad c_k = \frac{(-1)^k i}{k} \quad (k \neq 0)$$

Th.  $\phi_k(x) = e^{ikx}$  is complete in  $L^2([-pi; pi])$   $k \in I \equiv \mathbb{Z}$  if

$$\sum_{\substack{k=-\infty \\ k \neq 0}}^{+\infty} \left| \frac{(-1)^k i}{k} \right|^2 \cdot \underbrace{\|\phi_k\|_2^2}_{2\pi} = \|f\|_2^2$$

$\downarrow \int_{-\pi}^{\pi} x^2 dx$

$$2\pi \sum_{k \neq 0} \frac{1}{k^2} = \left[ \frac{x^3}{3} \right]_{-\pi}^{\pi}$$

$$2\pi \left( \sum_{k=1}^{\infty} \frac{1}{k^2} + \sum_{k=-1}^{-\infty} \frac{1}{k^2} \right) = \frac{1}{3} (\pi^3 + \pi^3)$$

$$4\pi \sum_{k=1}^{\infty} \frac{1}{k^2} = \frac{2}{3} \pi^3$$

$\frac{\pi^2}{6}$

$$\frac{4\pi^3}{6} = \frac{2}{3} \pi^3$$

The equality is true,  
so the orthogonal system is complete!

(def.) (1) (2) (3)  
Complete Orthogonal Systems of functions in  
 $L^2$ ,  $\phi_\kappa(x)$

We say that  $\{\phi_\kappa\}$ ,  $\kappa \in I$ , is ortho-normal if

$$\|\phi_\kappa\|_2 = 1 \quad \forall \kappa \in I$$

N.B. Any orthogonal system can be normalized  
 just multiplying  $\phi_\kappa(x)$  by a constant  $\frac{1}{\|\phi_\kappa\|_2}$

Th. 2.3.1 (p. 33) (minimizing property of Fourier coeff's)

Suppose  $\{\phi_\kappa(x)\}_{\kappa \in I}$ ,  $x \in (a, b)$ , is an orthogonal

system,  $f \in L^2(a, b)$ ,  $f \sim \sum_{\kappa \in I} c_\kappa \phi_\kappa(x)$

$$\Rightarrow \|f(x) - \sum_{\kappa=-n}^n c_\kappa \phi_\kappa(x)\|_2 \stackrel{(*)}{\leq} \|f(x) - \sum_{\kappa=-n}^n d_\kappa \phi_\kappa(x)\|_2$$

and the = in (\*) holds  $\Leftrightarrow c_\kappa = d_\kappa \quad \forall \kappa$

N.B.  $c_\kappa$  is the (generalized) Fourier coefficient of  $f(x)$  with respect to the orthogonal system  $\{\phi_\kappa(x)\}$  i.e.

$$c_\kappa = \frac{1}{\|\phi_\kappa\|_2^2} \cdot (f, \phi_\kappa)$$

(\*) means the "square deviation" of  $f(x)$  from

$\sum_{\kappa=-n}^n d_\kappa \phi_\kappa(x)$  is minimal when  $d_\kappa = c_\kappa$ , where

$c_\kappa$  are the Fourier coeff. of  $f$  w.r.t.  $\{\phi_\kappa\}$   
 (with respect to)

Proof:  $\|f - \sum_{k=-n}^n d_k \phi_k\|_2^2 = (f - \sum_{k=-n}^n d_k \phi_k, f - \sum_{k=-n}^n d_k \phi_k) =$   
 $\|f(x) - \sum_{k=-n}^n d_k \phi_k(x)\|_2^2$  (scalar product)

using distributive property

$$= (f, f) - (f, \sum_{k=-n}^n d_k \phi_k) - (\sum_{k=-n}^n d_k \phi_k, f) +$$

$$+ (\sum_{k=-n}^n d_k \phi_k, \sum_{k=-n}^n d_k \phi_k) = \|f\|_2^2 - \sum_{k=-n}^n (f, d_k \phi_k) +$$

$$- \sum_{k=-n}^n (d_k \phi_k, f) + \sum_{j=-n}^n \sum_{k=-n}^n (d_k \phi_k, d_j \phi_j) =$$

Remark

$$(\alpha f, g) = \alpha (f, g)$$

$$(f, \alpha g) = \bar{\alpha} (f, g)$$

$$= \|f\|_2^2 - \sum_{k=-n}^n \bar{d}_k (f, \phi_k) - \sum_{k=-n}^n d_k (\phi_k, f) +$$

$$+ \sum_{k=-n}^n \sum_{j=-n}^n d_k \bar{d}_j (\phi_k, \phi_j) = (*) \quad (\phi_k, \phi_j) \neq 0 \text{ when } k=j$$

Remark

$$(f, \phi_k) = c_k \|\phi_k\|_2^2 \quad (\text{Fourier coeff. } c_k \text{ from definition})$$

$$(\phi_k, f) = \overline{(f, \phi_k)} = \bar{c}_k \|\phi_k\|_2^2$$

$$(*) = \|f\|_2^2 - \sum_{k=-n}^n \bar{d}_k c_k \|\phi_k\|_2^2 - \sum_{k=-n}^n d_k \bar{c}_k \|\phi_k\|_2^2 +$$

$$+ \sum_{k=-n}^n d_k \bar{d}_k \|\phi_k\|_2^2 = + \sum_{k=-n}^n c_k \bar{c}_k \|\phi_k\|_2^2 - \sum_{k=-n}^n c_k \bar{d}_k \|\phi_k\|_2^2$$

collect these terms

$$= \|f\|_2^2 - \sum_{k=-n}^n c_k \bar{c}_k \|\phi_k\|_2^2 + \sum_{k=-n}^n (c_k - d_k) (\bar{c}_k - \bar{d}_k) \|\phi_k\|_2^2 =$$

$$= \|f\|_2^2 - \sum_{k=-n}^n |c_k|^2 \|\phi_k\|_2^2 + \sum_{k=-n}^n |c_k - d_k|^2 \|\phi_k\|_2^2 \quad (\star)$$

This quantity is minimal  $\iff c_k = d_k$  for  $k = -n, -n+1, \dots, n-1, n$

(Q.E.D.)

Furthermore If  $c_k = d_k$   $S_m f(x) = \sum_{k=-m}^m c_k \phi_k(x)$

Partial Sums Order  $m$  (bilateral)  
of our F.S. (generalized)

Then the last term in (\*) formula is  $= 0 \Rightarrow$

$$\|f(x) - S_m f(x)\|_2^2 = \|f\|_2^2 - \sum_{k=-m}^m |c_k|^2 \|\phi_k\|_2^2 \quad (\text{proof of corollary})$$

This observation ("furthermore") implies proof of Bessel's

Inequality:  $\|f\|_2^2 \geq \sum_{k=-m}^m |c_k|^2 \|\phi_k\|_2^2$

Def. we say that  $\{\phi_k\}_{k \in I}$  has the Parseval Property

if  $=$  holds in  $\|f\|_2^2 = \sum_{k \in I} |c_k|^2 \|\phi_k\|_2^2$

Theorem (no proof in this 5 credits course!)

An orthogonal system of functions  $\{\phi_k(x)\}_{k \in I}$  is complete  $\Leftrightarrow$  it has the Parseval Property.

Let's revisit these facts in the classical case when  $\{\phi_k(x) = e^{ikx}\}$  and  $I = \mathbb{Z} = \{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\}$

Also, we observed that, in the classical case, we can write down in 2 forms the F.S. of  $f(x)$

$$\sim \sum_{k=-\infty}^{+\infty} c_k e^{ikx} = \frac{a_0}{2} + \sum_{k=1}^{\infty} (a_k \cos kx + b_k \sin kx)$$

N.B.  $\{\cos kx, \sin kx, \frac{1}{2}\}_{k=1,2,3,\dots}$  is a complete orthogonal system in  $L^2(-\pi; \pi)$ . And, of course, also  $\{e^{ikx}\}_{k \in \mathbb{Z}}$ .

Parseval  $\|f\|_2^2 = \int_{-\pi}^{\pi} |f(x)|^2 dx = \sum_{k=-\infty}^{+\infty} |c_k|^2 \|\phi_k\|_2^2$   
 $\{e^{ikx}\}$   $\underbrace{\qquad\qquad\qquad}_{2\pi \forall k}$

i.e.  $\frac{1}{2\pi} \int_{-\pi}^{\pi} |f(x)|^2 dx = \sum_{n=-\infty}^{+\infty} |c_n|^2$

Parseval  $\left\{ \cos nx, \sin nx, \frac{1}{2} \right\}$

We proved that

$$c_n = \frac{a_n - ib_n}{2} \quad \text{for } n > 0 \quad (c_0 = a_0/2) \quad b_0 = 0$$

$$c_n = \frac{a_n + ib_n}{2} \quad \text{for } n < 0$$

$$a_n = c_n + c_{-n} \quad \text{for } n > 0 \quad (a_0 = 2c_0)$$

$$b_n = i(c_n - c_{-n}) \quad \text{for } n > 0 \quad (b_0 = 0)$$

$$\sum_{n=-\infty}^{+\infty} |c_n|^2 = \sum_{n=1}^{\infty} |c_n|^2 + |c_0|^2 + \sum_{n=1}^{\infty} |c_{-n}|^2 = \text{we substitute the previous relationships} =$$

$$= \sum_{n=1}^{\infty} \frac{|a_n + ib_n|^2}{4} + \frac{|a_0|^2}{4} + \sum_{n=1}^{\infty} \frac{|a_n - ib_n|^2}{4} =$$

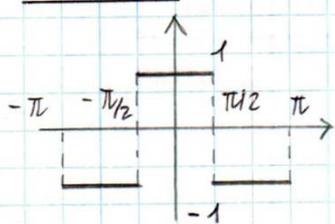
by def. of complex number modulus;  
 $z = x + iy$   
 $|z| = \sqrt{x^2 + y^2}$

$$= \frac{|a_0|^2}{4} + \frac{1}{2} \sum (|a_n|^2 + |b_n|^2), \text{ because } \frac{1}{4} [(a_n + ib_n)(\bar{a}_n - i\bar{b}_n) + (a_n - ib_n)(\bar{a}_n + i\bar{b}_n)] =$$

$$= \frac{1}{4} [ |a_n|^2 - i a_n \bar{b}_n + i b_n \bar{a}_n + |b_n|^2 + |a_n|^2 + i a_n \bar{b}_n - i b_n \bar{a}_n + |b_n|^2 ] = \frac{1}{2} (|a_n|^2 + |b_n|^2)$$

(Real) Parseval  $\frac{1}{2\pi} \int_{-\pi}^{\pi} |f(x)|^2 dx = \frac{|a_0|^2}{4} + \frac{1}{2} \sum_{n=1}^{\infty} (|a_n|^2 + |b_n|^2)$

Example



square wave

F.S.  $f(x) = \frac{4}{\pi} \left( \cos x - \frac{\cos 3x}{3} + \frac{\cos 5x}{5} - \dots \right)$

Parseval  $\frac{1}{2\pi} \int_{-\pi}^{\pi} |f(x)|^2 dx =$   
 $(b_n = 0)$   
 $(a_0 = 0)$

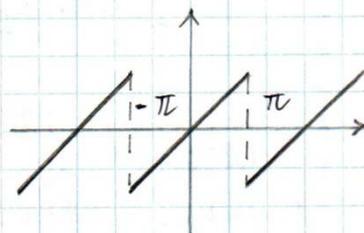
$$= \frac{1}{2} \frac{16}{\pi^2} \left( \frac{1}{1^2} + \frac{1}{3^2} + \frac{1}{5^2} + \frac{1}{7^2} + \dots \right)$$

where  $a_n = \frac{(-1)^{n+1}}{n} \cdot \frac{4}{\pi}$

$$\underbrace{\frac{1}{2\pi} \int_{-\pi}^{\pi} 1 dx}_{=1} \Rightarrow \frac{1}{1^2} + \frac{1}{3^2} + \frac{1}{5^2} + \dots = \frac{\pi^2}{8}$$

Other ex. sawtooth

$$f(x) = 2 \left( \sin x - \frac{\sin 2x}{2} + \frac{\sin 3x}{3} - \dots \right)$$



$$\frac{1}{2\pi} \int_{-\pi}^{\pi} x^2 dx = \frac{1}{2} \sum_{n=1}^{\infty} 2^2 \frac{1}{n^2}$$

$$\frac{1}{2\pi} \left[ \frac{x^3}{3} \right]_{-\pi}^{\pi} = \frac{4}{2} \left( \frac{1}{1^2} + \frac{1}{2^2} + \frac{1}{3^2} + \dots \right)$$

$$\frac{1}{4\pi} \left( \frac{\pi^3 - (-\pi)^3}{3} \right) = \frac{1}{1^2} + \frac{1}{2^2} + \frac{1}{3^2} + \dots$$

$$\Rightarrow \frac{\pi^2}{6} = \frac{1}{1^2} + \frac{1}{2^2} + \frac{1}{3^2} + \dots$$

Theorem Schwarz Inequality (Cauchy Schwarz Buniaowski)

$$|(f, g)| \leq \|f\|_2 \|g\|_2 \quad (\text{scalar product})$$

$$\text{i.e.} \quad \left| \int_a^b f(x) \overline{g(x)} dx \right| \leq \left( \int_a^b |f(x)|^2 dx \right)^{1/2} \left( \int_a^b |g(x)|^2 dx \right)^{1/2}$$

N.B. this inequality is "obvious" if we take for granted the euclidean ( $\infty$  dimensional) nature of scalar products and  $L^2$  norms, otherwise not so obvious...

Remark The Schwarz Inequality is also true if our scalar product contains a weight function  $w(x) > 0$

$$\text{i.e.} \quad (f, g) \stackrel{\text{def.}}{=} \int_a^b f(x) \overline{g(x)} w(x) dx$$

in particular  $w(x) = 1$

$$\left| \int_a^b f(x) \overline{g(x)} w(x) dx \right| \leq \left( \int_a^b |f(x)|^2 w(x) dx \right)^{1/2} \left( \int_a^b |g(x)|^2 w(x) dx \right)^{1/2}$$

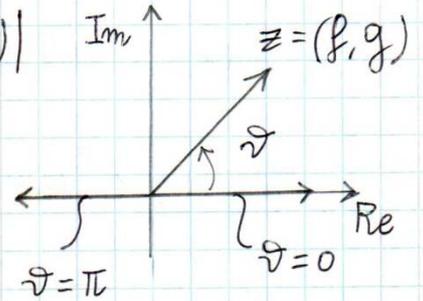
Proof: set  $\vartheta = \text{Arg}(f, g)$

[N.B. if  $(f, g)$  is real  $\vartheta = 0$  if  $(f, g) > 0$   
 $\vartheta = \pi$  if  $(f, g) < 0$ ]

Then  $e^{-i\vartheta} (f, g) = (e^{-i\vartheta} f, g) \equiv |(f, g)|$

Produces a rotation  $\vartheta$

coincides with the abs. val.  
 in this case  $\vartheta = 0; \pi$



$$z_1 = \int_1 e^{i\vartheta_1}$$

$$z_2 = \int_2 e^{i\vartheta_2}$$

$$z_1 z_2 = \int_1 \int_2 e^{i(\vartheta_1 + \vartheta_2)}$$

DE MOIVRE FORMULA

Now,  $e^{-i\vartheta}$  is a complex number of argument  $-\vartheta$  and  $|e^{-i\vartheta}| = 1$

So the effect of multiplying  $z \in \mathbb{C}$  by  $e^{-i\vartheta}$  is to rotate the vector  $z$  (in the complex plane) of an angle  $-\vartheta$

If  $\text{Arg } z = \vartheta$ , then  $e^{-i\vartheta} z = |z|$

Also, in this case, ( $\vartheta = 0; \pi$ )

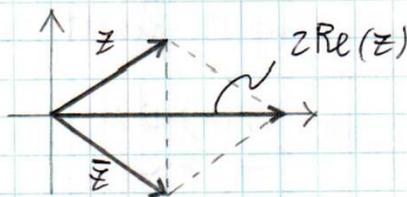
$$(f e^{-i\vartheta}, g) = \text{Re}(f e^{-i\vartheta}, g) = |(f, g)|$$

Now we go on with our proof:

$$0 \leq \| f e^{-i\vartheta} + \lambda g \|_2^2$$

(Where  $\lambda \in \mathbb{R}$   
 $\vartheta = \text{Arg}(f, g)$ )

$$\begin{aligned} &= (f e^{-i\vartheta} + \lambda g, f e^{-i\vartheta} + \lambda g) = \text{distributive property} = \\ &= (f e^{-i\vartheta}, f e^{-i\vartheta}) + \lambda (f e^{-i\vartheta}, g) + \lambda (g, f e^{-i\vartheta}) + \lambda^2 (g, g) = \\ &= \| f e^{-i\vartheta} \|_2^2 + 2\lambda \text{Re}(f e^{-i\vartheta}, g) + \lambda^2 \| g \|_2^2 = \end{aligned}$$



$$= \|f\|_2^2 + 2\lambda |(f, g)| + \lambda^2 \|g\|_2^2$$

$0 \leq$  polynomial of degree 2 in the variable  $\lambda$

(N.B. coeff. of  $\lambda^2$  is  $\|g\|_2^2 > 0$ )  $\Rightarrow \Delta \leq 0$

$$\frac{\Delta}{4} = \frac{b^2 - 4ac}{4} = \left(\frac{b}{2}\right)^2 - ac$$

$$\Rightarrow |(f, g)|^2 - \|f\|_2^2 \|g\|_2^2 \leq 0$$

$$|(f, g)|^2 \leq \|f\|_2^2 \|g\|_2^2$$

$$|(f, g)| \leq \|f\|_2 \|g\|_2 \quad (\text{Q.E.D.})$$

$$|(f, g)| \leq \|f\|_2 \|g\|_2$$

True also on  $L^2(\mathbb{R})$  or  $L^2([0, +\infty))$ ,  
so this inequality is relevant also  
for the Fourier Transform and  
Laplace Transform.

Lemma 2.4.2

(P.36)

$$(f, g) = \int_a^b f(x) \overline{g(x)} w(x) dx = \sum_{k \in I} c_k \overline{d_k} \| \phi_k \|_2^2$$

where  $c_k$  and  $d_k$  are the Fourier coeff. of  $f$  and  $g$   
with respect to the orthogonal system  $\{ \phi_k(x) \}$

$$f \sim \sum_{k \in I} c_k \phi_k(x) \quad \text{with } c_k = \frac{1}{\| \phi_k \|_2^2} (f, \phi_k)$$

$$\text{and } g \sim \sum_{k \in I} d_k \phi_k(x)$$

$$d_k = \dots$$

This formula (Parseval-Plancherel) in the special case  $f=g$

gives 
$$\int_a^b |f(x)|^2 dx = \sum_{k \in I} |c_k|^2 \|\phi_k\|_2^2$$

Proof

$| (f - S_n f, g) | \leq \|f - S_n f\|_2 \|g\|_2 \rightarrow 0$

we start writing this SCHWARZ INEQUALITY (whose general form we have proven before) as  $n \rightarrow \infty$

$$S_n f = \sum_{k=-n}^n c_k \phi_k(x)$$

(because  $\|g\|_2 > 0$  and  $\lim_{n \rightarrow \infty} \|f - S_n f\|_2 = 0$ )

$(f, g) = (f - S_n f + S_n f, g) = (f - S_n f, g) + (S_n f, g)$

therefore

$(f, g) = \lim_{n \rightarrow \infty} (S_n f, g) = \lim_{n \rightarrow \infty} \left( \sum_{k=-n}^n c_k \phi_k(x), g(x) \right) =$

$$= \lim_{n \rightarrow \infty} \sum_{k=-n}^n c_k \underbrace{(\phi_k, g)}_{\substack{|| \\ (g, \phi_k)}} = \lim_{n \rightarrow \infty} \sum_{k=-n}^n c_k \bar{d}_k \underbrace{\|\phi_k\|_2^2}_{\substack{\text{FROM THE DEFINITION OF THE} \\ \text{GENERALIZED COEFFS}}} =$$

$$= \sum_{k=-\infty}^{+\infty} c_k \bar{d}_k \|\phi_k\|_2^2$$

(QED)

Th. (No proof, but easy consequence of

what we have done today) if  $f \in L^1(a, b)$

and its F.S. is  $f \sim \sum_{k \in I} c_k \phi_k(x)$  suppose  $[a, \beta] \subset [a, b]$

then

$$\int_a^\beta f(x) dx = \int_a^\beta \sum_{k \in I} c_k \phi_k dx = \sum_{k \in I} c_k \int_a^\beta \phi_k(x) dx$$

{ if  $[a, \beta] \subset [a, b]$  finite intervals and  $f \in L^1([a, b]) \supset L^2([a, b])$  }  
 { then it's ok to exchange  $\int_a^\beta \dots$  with F. Series }

Remark In applications to ODE and PDE we would

also like to use 
$$[f(x)]' = \left[ \sum_{k \in I} c_k \phi_k(x) \right]' = \sum_{k \in I} c_k \phi_k'(x)$$

In general this is more delicate than exchanging F.S. and integrals ... But it's OK if  $C_k \rightarrow 0$  "quickly enough" (equivalent to  $f(x)$  being "smooth enough").

Some families of orthogonal systems  $\{\phi_k(x)\}_{k \in \mathbb{I}}$  which are not the classical  $\{e^{ikx}\}_{k \in \mathbb{Z}}$  or the equivalent  $\{\cos kx, \sin kx, \frac{1}{2}\}_{k=1,2,3,\dots}$

### Orthogonal Polynomials (pg. 45-59)

Idea: on a finite interval  $[a; b]$  the power functions  $1, x, x^2, x^3, \dots$  are linearly independent, i.e.,

$$\sum_{k=0}^n a_k x^k = 0 \text{ for all } x \in [a; b]$$

$$\Leftrightarrow a_k = 0 \text{ for } k=0, 1, 2, \dots, n$$

This fact is somehow connected to the fact that it is possible to approximate large spaces of functions using polynomials. For example if  $f \in L^p([a; b]) \forall \epsilon > 0 \exists m = m_\epsilon$  and a polynomial  $p(x)$  of degree  $m$  such that

$$\|f(x) - p(x)\|_{L^p([a; b])} < \epsilon \quad \left\{ \begin{array}{l} \text{N.B. the } L^\infty \text{ case is} \\ \text{UNIFORM approximation} \end{array} \right\}$$

Given a scalar product on  $L^2([a; b], w)$

$$(f, g) \stackrel{\text{def}}{=} \int_a^b f(x) \overline{g(x)} w(x) dx$$

usually is not true that the power functions are orthogonal with respect to each other.

BUT it is possible to find a sequence of polynomials

$p_0(x), p_1(x), p_2(x), p_3(x), \dots$ , where each  $p_k(x)$

is a polynomial of (exact) degree  $n$  (the coeff. of  $x^n$  in  $p_n(x)$  is  $\neq 0$ ) such that  $(p_n(x), p_j(x)) = \int_a^b p_n(x) \overline{p_j(x)} w(x) dx =$

$$= \begin{cases} 0 & \text{if } n \neq j \\ \alpha_{n>0} & \text{if } n = j \end{cases}$$

Furthermore if we fix  $[a; b]$  and a weight  $w(x) > 0$  for  $x \in [a; b]$  (possibly  $w(x) = 1$ ), the sequence  $p_0(x), p_1(x), p_2(x), \dots$  is essentially UNIQUE.

"essentially" in the sense that each  $p_n(x)$  can be dilated changing it into  $\beta_n p_n(x)$

If we normalize, e.g. choosing  $\|p_n\|_2 = 1 \quad \forall n = 0, 1, 2, \dots$

then the sequence is unique. (other normalizations are possible, e.g.,  $p_n(0) = 1$ )

Problem Given  $[a; b]$  and  $w(x)$ , how do we compute  $p_0(x), p_1(x), p_2(x), \dots$  orthogonal polynomials with respect to  $(f, g) = \int_a^b f(x) \overline{g(x)} w(x) dx$  ?

Several answers:

1<sup>st</sup>: "Bulldozer" answer: Gram-Schmidt Orthogonalization

Advantage: it "always" works (for any given  $[a; b] w(x)$ ) starting from the non-orthogonal but independent system  $1, x, x^2, x^3, \dots$

Disadvantage: boring, inefficient, slow (even on computers)

2<sup>nd</sup> possibility: It can be proven that all orthogonal polynomials satisfy a  $\geq 3$  terms recursive formula:

$$p_m \stackrel{*}{=} A_m p_{m-1}^{(x)} + B_m p_{m-2}^{(x)} \quad \neq \quad \neq$$

Recipe: compute  $p_0(x)$  and  $p_1(x)$ . From  $p_0$  and  $p_1(x)$  compute  $p_2(x)$  using  $*$   
From  $p_1$  and  $p_2$  compute  $p_3$  etc. ...

3<sup>rd</sup> method: Direct Formulae given  $n$  produce  $p_m$

4<sup>th</sup> method: Semidirect Formulae (Rodriguez)

$$p_m(x) = \frac{d^m}{dx^m} (\text{something})$$

Other methods ...

Depending on  $[a; b]$  and  $w(x)$  the "smart" method could change.

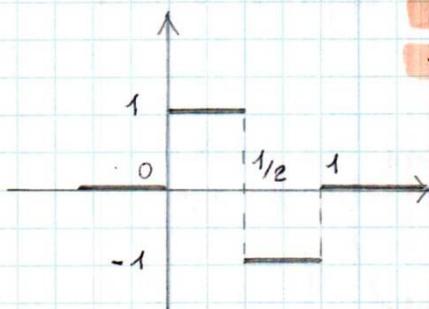
In this course we'll see only a small sample of these methods (mostly following the book)

Wavelets (in particular Haar Basis) [ALSO P.62]  
~1901

Theorem (Haar)

$$\psi(x) = \begin{cases} 1 & \text{if } x \in (0, \frac{1}{2}) \\ -1 & \text{if } x \in (\frac{1}{2}, 1) \\ 0 & \text{elsewhere} \end{cases}$$

MOTHER  
FUNCTION

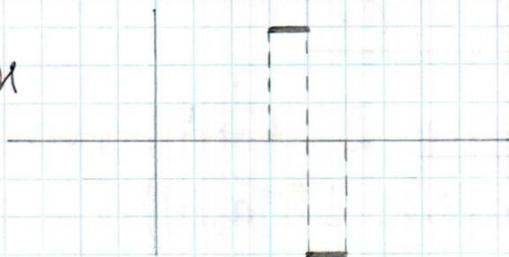


( $\psi$  is "one tooth" of a square wave)

Def.

$$\psi_{j\kappa}(x) = 2^{j/2} \psi(2^j x - \kappa) \text{ for } j \in \mathbb{Z} \quad \kappa \in \mathbb{Z}$$

$\psi_{j\kappa}$



$\psi_{j\kappa}(x)$  is a complete orthonormal system in  $L^2(\mathbb{R})$

in other words any  $f \in L^2(\mathbb{R})$  can be written

$$f(x) = \sum_{j,k} c_{jk} \psi_{jk}(x), \text{ where the "Haar coefficients"}$$

(Generalized F.S.)

$$c_{jk} = (f, \psi_{jk}) =$$

$$= \int_{-\infty}^{+\infty} f(x) \overline{\psi_{jk}(x)} dx = \int f(x) \psi_{jk}(x) dx$$

(real function)      interval support of  $\psi_{jk}$

(N.B.  $\|\psi_{jk}\|_2^2 = 1$   
because  $\{\psi_{jk}\}$  is  
an orthonormal system)

Advantage of Haar system: very simple  $\psi(x)$

Disadvantages: many coefficients are needed to get a good approximation of  $f(x)$ . Properties of  $f(x)$  (like smoothness) are mapped into  $c_{jk}$  in a complicated way.

Modern approach: choose better  $\psi(x)$  mother functions such that  $\psi_{jk}(x) = 2^{j/2} \psi(2^j x - k)$  for  $j, k \in \mathbb{Z}$  is a complete orthonormal system in  $L^2(\mathbb{R})$ .

Other big families of orthogonal systems are given by eigenfunctions of specific operators.

COMPARE  
PAGE  
41R &  
44

Def. If  $A$  is an  $n \times n$  matrix it defines linear operator:  $\mathbb{R}^n \rightarrow \mathbb{R}^n$

$$\begin{pmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{pmatrix} = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}$$

$\alpha$  is an EIGENVALUE and  $\vec{x}$  is an EIGENVECTOR of this operator if  $\boxed{A\vec{x} = \alpha\vec{x}}$

Def. Given a linear operator  $L: A \rightarrow A$ , where  $A$  is a space of functions (e.g.  $L^p([a, b])$ ...)

Linear  $\mathcal{L}(\alpha f(x) + \beta g(x)) = \alpha \mathcal{L}(f(x)) + \beta \mathcal{L}(g(x))$

(ex. of linear operators: derivative, n-th derivative, Fourier transform, Laplace Transform ...)

If  $\mathcal{L}(f(x)) = \alpha \cdot f(x)$ , then we say that  $\alpha$  is an EIGENVALUE and  $f$  is an EIGENFUNCTION of  $\mathcal{L}$ .

ex.  $f(x) = e^{ax}$   $\mathcal{L} = \text{derivative}$

$$\mathcal{L}(e^{ax}) = a e^{ax} \quad \text{eigenvalue} = a$$

$\mathcal{L} = \text{second derivative}$   $f(x) = \cos x, \sin x$   
eigenfunctions

$$(\sin x)'' = -\sin x \quad \mathcal{L}(\sin x) = (-1) \sin x$$

$$\text{eigenvalue} = -1$$

$$\mathcal{L}(\cos x) = -\cos x \quad \text{eigenvalue} = -1$$

N.B.  $\mathcal{L}(3 \cos x) = 3(-\cos x) = (-1)(3 \cos x)$

↑ same eigenvalue

ex.  $\mathcal{L} = \text{Fourier Transform}$

$$\mathcal{L}(e^{-ax^2}) = \sqrt{\frac{\pi}{a}} e^{-\frac{t^2}{4a}} \quad (\text{Gaussian function})$$

$a > 0$  let's choose  $a = \frac{1}{2}$

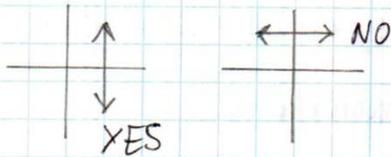
$$\mathcal{L}(e^{-\frac{1}{2}x^2}) = \sqrt{2\pi} e^{-\frac{1}{2}t^2}$$

so  $f(x) = e^{-x^2/2}$  eigenfunction eigenvalue =  $\sqrt{2\pi}$

Rem. 1: we have to choose  $a = \frac{1}{2}$  to have  $f(x)$  eigenfunction

of the F.T.; the corresponding eigenvalue is  $\sqrt{2\pi}$  BUT  $\Rightarrow$

$$\mathcal{L}(c e^{-\frac{1}{2}x^2}) = c \sqrt{2\pi} e^{-\frac{1}{2}\omega^2}$$



Rem. 2: (more details later) it is possible to find polynomials  $h_0(x), h_1(x), h_2(x), \dots, h_n(x)$  of degree  $n$  such that

P.43R

$$\text{F.T.}(h_n(x) e^{-x^2/2}) = a_n \underbrace{h_n(x)}_{\text{HERMITE POLYNOMIALS}} e^{-x^2/2} \quad \text{HERMITE FUNCTIONS}$$

It is possible to prove also that these eigenfunctions of the F.T.  $h_n(x) e^{-x^2/2}$  are a complete orthogonal system in  $L^2(\mathbb{R})$ .



Orthogonal polynomials

$p_n(x) = c_n x^n + \dots + c_1 x + c_0$  (linear combination of powers degree =  $n$  = highest power)

Given  $[a; b]$  (also  $[0; +\infty)$ , also  $(-\infty; +\infty)$ ) and a weight (i.e. a function  $w(x) \geq 0$  for  $x \in [a; b]$ ) we have a scalar product  $(f, g) = \int_a^b f(x) \overline{g(x)} w(x) dx$ ; we also have an "essentially" unique sequence of polynomials of degree  $n = 0, 1, 2, 3, \dots$  such that it is an orthogonal system,

$$\text{i.e. } (p_m, p_n) = \int_a^b p_m(x) \overline{p_n(x)} w(x) dx = \begin{cases} 0 & \text{if } m \neq n \\ \|p_m\|_2^2 & \text{if } m = n \end{cases}$$

Depending on the weight choice we get different solutions.

case n. 1 Legendre Polynomials:  $[a; b] = [-1; 1]$   $w(x) = 1 \quad \forall x \in [-1; 1]$   
our scalar product is  $(f, g) = \int_{-1}^1 f(x) \overline{g(x)} dx$ . They are:

$$p_0(x) = 1; \quad p_1(x) = x; \quad p_2(x) = \frac{3}{2}x^2 - \frac{1}{2}; \quad p_3(x) = \frac{5}{2}x^3 - \frac{3}{2}x;$$

$$p_4(x) = \frac{35}{8}x^4 - \frac{15}{4}x^2 + \frac{3}{8}$$

(are chosen in such a way that they make oscillation in  $[-1; 1]$ )

Let's check orthogonality in one case  $(p_2, p_3) = \int_{-1}^1 \left(\frac{3}{2}x^2 - \frac{1}{2}\right) \left(\frac{5}{2}x^3 - \frac{3}{2}x\right) dx$

$$= \int_{-1}^1 \frac{15}{4}x^5 - \frac{9}{4}x^3 - \frac{5}{4}x^3 + \frac{3}{4}x dx = \left[ \frac{15}{4} \frac{x^6}{6} - \frac{14}{4} \frac{x^4}{4} + \frac{3}{4} \frac{x^2}{2} \right]_{-1}^1 =$$

$$\underbrace{\frac{15}{4} \frac{x^6}{6} - \frac{14}{4} \frac{x^4}{4}}_{-\frac{14}{4}x^3} + \frac{3}{4} \frac{x^2}{2}$$

$$= \frac{5}{8} - \frac{7}{8} + \frac{3}{8} - (\text{same things}) = 0$$

$$(p_2, p_2) = \int_{-1}^1 \left( \frac{3}{2}x^2 - \frac{1}{2} \right)^2 dx > 0$$

These are polynomials, but, in the interval they act, they are imitating sine & cosine (oscillation, cancellation effect...)

See figure 3.1, page 47 of the book.

To obtain O.P. we could apply the Gram-Schmidt orthogonalization to the sequence of powers  $1, x, x^2, x^3, \dots$   
(N.B. these polynomials are linearly independent but they are not orthogonal w.r.t.  $(f, g) = \int_{-1}^1 f(x) \overline{g(x)} dx$ )

There are more efficient methods to obtain the same sequence of polynomials.

It can be proven that 
$$p_n(x) \stackrel{(*)}{=} \frac{1}{2^n \cdot n!} \frac{d^n}{dx^n} (x^2 - 1)^n$$
 (Rodriguez)

(hidden recursion...)

↳ because you can't compute directly the  $n$ -th derivative.

$p_0(x) = 1$  From the formula  $(*)$  (Rodriguez):

$$p_1(x) = \frac{1}{2^1 \cdot 1!} \frac{d}{dx} (x^2 - 1) = \frac{1}{2} 2x = x \quad \checkmark$$

$$\begin{aligned} p_2(x) &= \frac{1}{2^2 \cdot 2!} \frac{d^2}{dx^2} (x^2 - 1)^2 = \frac{1}{8} \frac{d^2}{dx^2} (x^4 - 2x^2 + 1) = \frac{1}{8} (4x^3 - 4x)' = \\ &= \frac{1}{8} (12x^2 - 4) = \frac{3}{2}x^2 - \frac{1}{2} \quad \checkmark \end{aligned}$$

Theorem (no proof in this course)

$$(\star) \quad \boxed{n p_n(x) - x(2n-1) p_{n-1}(x) + (n-1) p_{n-2}(x) = 0} \quad \left( \begin{array}{l} \text{3-terms} \\ \text{recursion} \end{array} \right)$$

Special case of Theorem  $\star$ : Almost all systems of O.P. satisfy a 3 terms recursion  $\left( \begin{array}{c} \circ \\ \parallel \\ \circ \end{array} \right)$

From  $(\star)$  if we know the first two cases  $p_0(x) = 1$ ,  $p_1(x) = x$  then recursively we can compute  $p_2(x)$ ,  $p_3(x)$ , ...

# Legendre Polynomials

Rodriguez  $P_n(x) = \frac{1}{2^n n!} \frac{d^n}{dx^n} (x^2-1)^n$

$$P_0(x) = 1$$

$$P_1(x) = \frac{1}{2} \frac{d}{dx} (x^2-1) = x$$

$$P_2(x) = \frac{1}{2^2 2!} \frac{d^2}{dx^2} (x^2-1)^2 = \frac{1}{8} \frac{d}{dx} [2(x^2-1) \cdot 2x] =$$
$$= \frac{1}{8} [4x^3 - 4x] = \frac{1}{8} [12x^2 - 4] = \frac{3}{2}x^2 - \frac{1}{2}$$

$$P_3(x) = \frac{1}{2^3 3!} \frac{d^3}{dx^3} (x^2-1)^3 = \frac{1}{8 \cdot 6} \frac{d^2}{dx^2} [3(x^2-1)^2 \cdot 2x] =$$
$$(x^4 + 1 - 2x^2) 6x$$
$$(6x^5 + 6x - 12x^3)$$
$$= \frac{1}{48} \frac{d}{dx} (30x^4 + 6 - 36x^2) =$$

$$= \frac{1}{48} (120x^3 - 72x) = \frac{5}{2}x^3 - \frac{3}{2}x$$

$$P_4(x) = \frac{1}{2^4 4!} \frac{d^4}{dx^4} (x^2-1)^4 = \frac{1}{384} \frac{d^3}{dx^3} [8x(x^2-1)^3] =$$

$$= \frac{1}{384} \frac{d^2}{dx^2} [8(x^2-1)^3 + 8x \cdot 3(x^2-1)^2 \cdot 2x] =$$
$$48x^2(x^2-1)^2$$

$$= \frac{1}{384} \frac{d}{dx} [24(x^2-1)^2 \cdot 2x + 96x(x^2-1)^2 + 48x^2 \cdot 2(x^2-1) \cdot 2x] =$$
$$144x(x^2-1)^2 + 192x^3(x^2-1)$$

$$= \frac{1}{384} \cdot [144(x^2-1)^2 + 144x \cdot 2(x^2-1)2x + 576x^2(x^2-1) +$$
$$+ 192x^3 \cdot 2x] =$$

$$= \frac{1}{384} [144x^4 + 144 - 288x^2 + 576x^4 - 576x^2 + 576x^4 - 576x^2 + 384x^4]$$

...  $\frac{144x^4}{144 \cdot 4} + \frac{144x^0}{144 \cdot 2} - \frac{288x^2}{144 \cdot 2} - \frac{35x^4}{15 \cdot 4} - \frac{15x^2}{15 \cdot 2} + \frac{3x^0}{2}$

### 3 term recursion

$$n p_n(x) - x(2n-1) p_{n-1}(x) + (n-1) p_{n-2}(x) = 0$$

$$p_0(x) = 1 \quad p_1(x) = x$$

---

$$2 p_2 - x(2 \cdot 2 - 1) x + (2 - 1) 1 = 0$$

$$2 p_2 = 3x^2 - 1 \quad p_2 = \frac{3x^2}{2} - \frac{1}{2}$$

---

$$3 p_3 - x(6-1) \left( \frac{3}{2} x^2 - \frac{1}{2} \right) + (3-1) x = 0$$

$$3 p_3 = \frac{15}{2} x^3 - \frac{5}{2} x - 2x = \frac{15}{2} x^3 - \frac{9}{2} x$$

$$p_3 = \frac{5}{2} x^3 - \frac{3}{2} x$$

---

$$4 p_4 - x(8-1) \left( \frac{5}{2} x^3 - \frac{3}{2} x \right) + 3 \left( \frac{3x^2}{2} - \frac{1}{2} \right) = 0$$

$$4 p_4 - \frac{35}{2} x^4 + \frac{21}{2} x^2 + \frac{9}{2} x^2 - \frac{3}{2} = 0$$

$$4 p_4 = \frac{35}{2} x^4 - 15x^2 + \frac{3}{2}$$

$$p_4 = \frac{35}{8} x^4 - \frac{15}{4} x^2 + \frac{3}{8}$$

HIGHLY MORE EFFECTIVE PROCEDURE

Exercise: compute the first 5 or 6 (or 7) Legendre Polynomials both with Rodriguez Formula and the (\*) recursion comparing efficiency.

A little detour about recursive method: for example, Fibonacci Numbers  $C_0, C_1, C_2, C_3, \dots$

$$C_0 = C_1 = 1 \text{ and } C_n = C_{n-1} + C_{n-2}$$

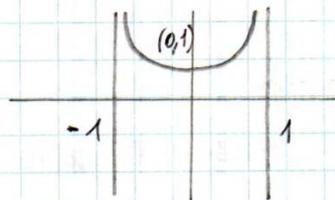
$$C_2 = 1+1=2 \quad C_3 = 2+1=3 \quad C_4 = \dots$$

exercise: either on your own (better) or cheating (Google) find a direct closed formula for the  $n$ -th Fibonacci Number  $C_n$  (hint: it is the linear combination of two exponentials one increasing to  $\rightarrow \infty$  and the other one decreasing to 0) (hint: Golden Ratio)

Sometimes (but not always) it is possible, like in the case of Fibonacci, to obtain a "closed formula" from the recursive definition.

Case 2: Chebyshev Polynomials  $T_n(x)$  orthogonal system of polynomials in  $(-1; 1)$  with the weight

$$W(x) = \frac{1}{\sqrt{1-x^2}}$$



$$(f, g) = \int_{-1}^1 f(x) \overline{g(x)} \frac{1}{\sqrt{1-x^2}} dx \quad (**)$$

$$T_0(x) = 1$$

$$T_1(x) = x$$

$$T_2(x) = 2x^2 - 1$$

$$T_3(x) = 4x^3 - 3x$$

$$T_4(x) = 8x^4 - 8x^2 + 1$$

$$T_5(x) = \dots$$

First method to compute these polynomials is Gram-Schmidt apply to  $1, x, x^2, x^3, \dots$  w.r.t. (\*\*)

2nd method

$$T_{m+1}(x) - 2xT_m(x) + T_{m-1}(x) = 0$$

(try it)

$$T_2(x) - 2xT_1(x) + T_0(x) = 0$$

Theorem "closed formula" for Chebycheff Polynomials

$$T_m(x) = \cos(m \arccos x) = \dots =$$

$$= x^m - \binom{m}{2} x^{m-2} (1-x^2) + \binom{m}{4} x^{m-4} (1-x^2)^2 - \dots$$

Case 3: Hermite Polynomials (and Hermite Functions)

consider the interval  $(-\infty; +\infty) \equiv \mathbb{R}$   $w(x) = e^{-x^2}$



$$(f, g) = \int_{-\infty}^{+\infty} f(x) \overline{g(x)} e^{-x^2} dx$$

we get Hermite Polynomials  $H_m(x)$

$$H_{m+1}(x) - 2xH_m(x) + 2mH_{m-1}(x) = 0$$

$$H_0(x) = 1 \quad H_1(x) = 2x$$

Rodriguez: 
$$H_m(x) = (-1)^m e^{x^2} \frac{d^m}{dx^m} e^{-x^2}$$

e.g. 
$$H_1(x) = (-1)^1 e^{x^2} \frac{d}{dx} (e^{-x^2}) = -e^{x^2} (-2x) e^{-x^2} = 2x$$

$$H_2(x) = (-1)^2 e^{x^2} \frac{d}{dx} (-2xe^{-x^2}) = e^{x^2} (-2e^{-x^2} - 2x(-2x)e^{-x^2})$$
$$= -2 + 4x^2$$

Important Remark Once we have generated the sequence

$H_0(x), H_1(x), H_2(x), \dots$  of Hermite Polynomials we know that

$$\int_{-\infty}^{+\infty} H_m(x) H_n(x) e^{-x^2} dx = \begin{cases} 0 & n \neq m \\ \|H_m\|_2^2 > 0 & \text{if } n = m \end{cases}$$

We can also consider the  $h_m(x) = H_m(x) e^{-x^2/2}$  called

Hermite Functions (N.B. these functions are not polynomials, but they are polynomials times the square root of a weight)

$$h_0(x) = e^{-x^2/2} ; h_1(x) = 2xe^{-x^2/2} ; h_2(x) = (4x^2 - 2)e^{-x^2/2}$$